

# *Timestamped Event Stream System*

## *Chronicle Recognition System*

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# Outline

## Architecture

- interface producer vs. consumer

## Feedback : control by consumers

- messages of the control protocol

## Chronicles

- focusing when computing correlation rules



# 1

## Architecture



# Producer/Consumer

## Define the way of event transmission

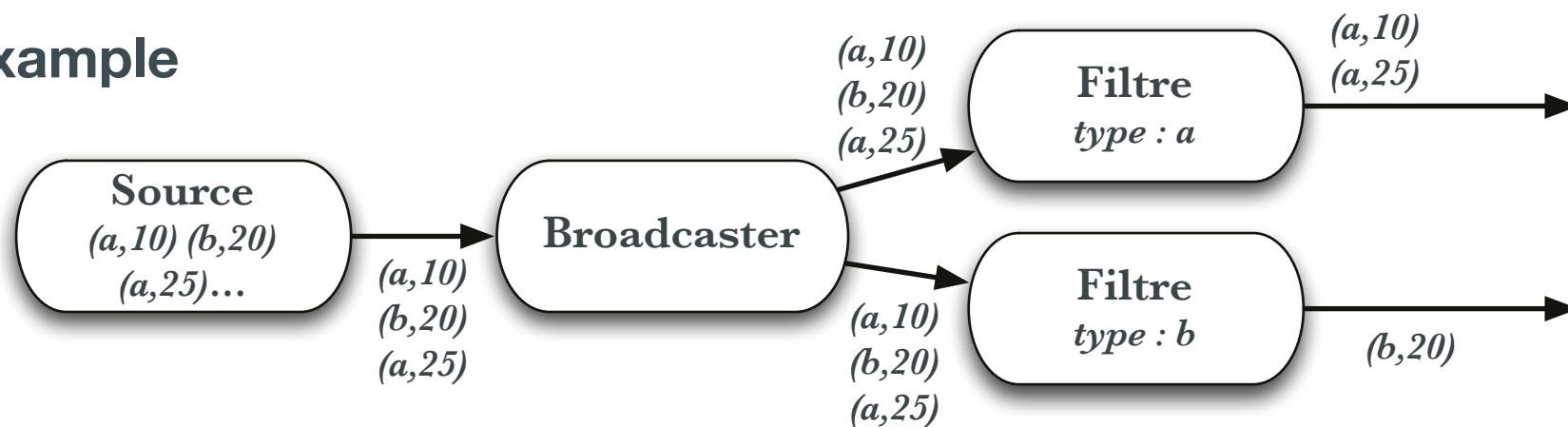


contract: an event  $(e,t)$  is sent only one time

- notice: distinct dates mean different events



## Example



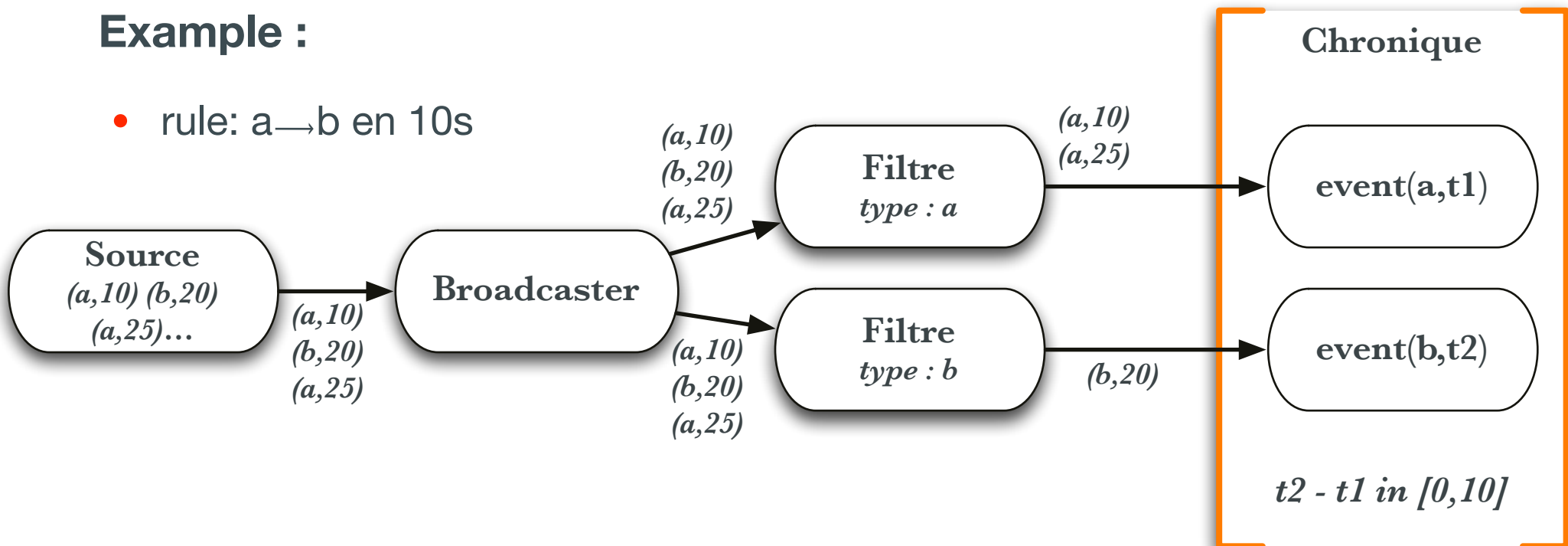
# Application to the correlation of alarms

The correlation rule processing is based on

- a set of consumers (one per event of the rule)
- correlation information (e.g. the time constraint)

Example :

- rule:  $a \rightarrow b$  en 10s

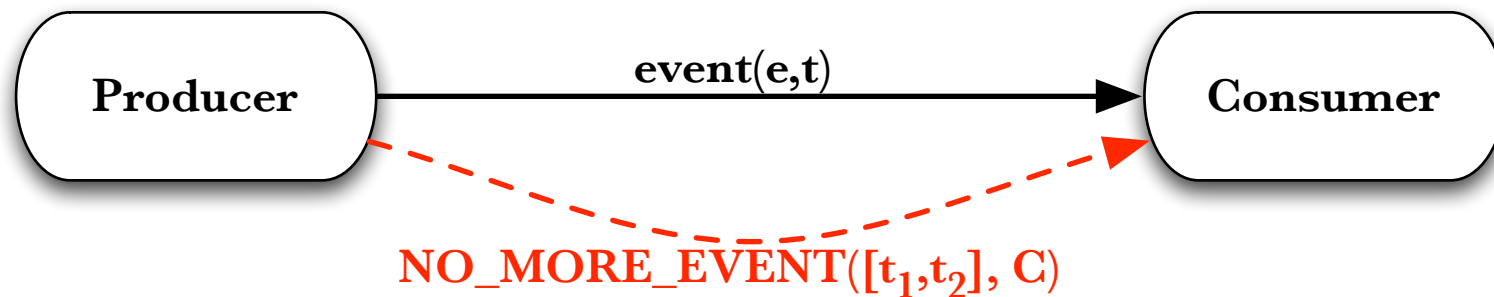


# Control/sync. of observations

## First control message : to sync. the streams

 *contract: the producer will never send an event which satisfies  $C$  and with an occurrence date within  $TW$*

- *This message could consider as a closure control message.*



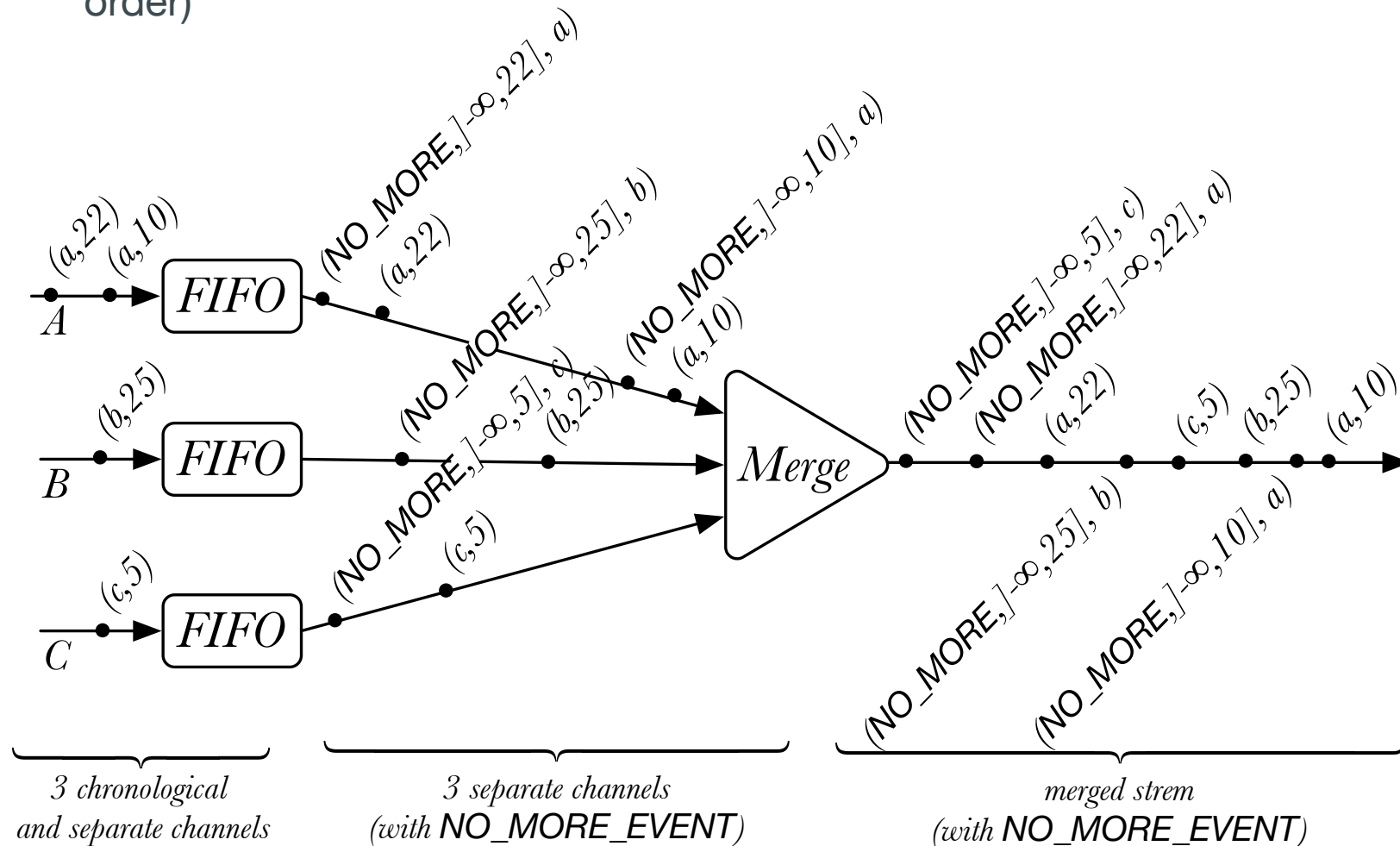
- **Advantage:** event transmission policy could be non chronological
  - `anm([0,10]); ... anm([100,200]);... anm([0,50])`



# NO\_MORE\_EVENT Usage (1/2): Multiple Sources

## Example of user : streams joint

- hyp. each stream is FIFO (i.e. each source sends events in a chronological order)



# NO\_MORE\_EVENT Usage (2/2): Deduced Events

## Example:


- chronicle C
  - $\text{event}(a,t) \wedge \text{event}(b,t') \wedge (t-t' \text{ in } [90,120]) \succ (c, t')$
- Receiving  $(a,10)$ 
  - a correlation context starts with  $t=10$  and  $t' \in [100,130]$
- Receiving  $\text{assertNoMore}([-\infty, 50])$ 
  - the previous context remains valid (and always waits for b)
  - **nevertheless we already be able to deduce  $\text{assertNoMore}(W, \text{type:c})$**   
with  $W = ([-\infty,50] \dot{+} 90) \setminus [100,130] = ]-\infty,100[ \cup ]130,140]$



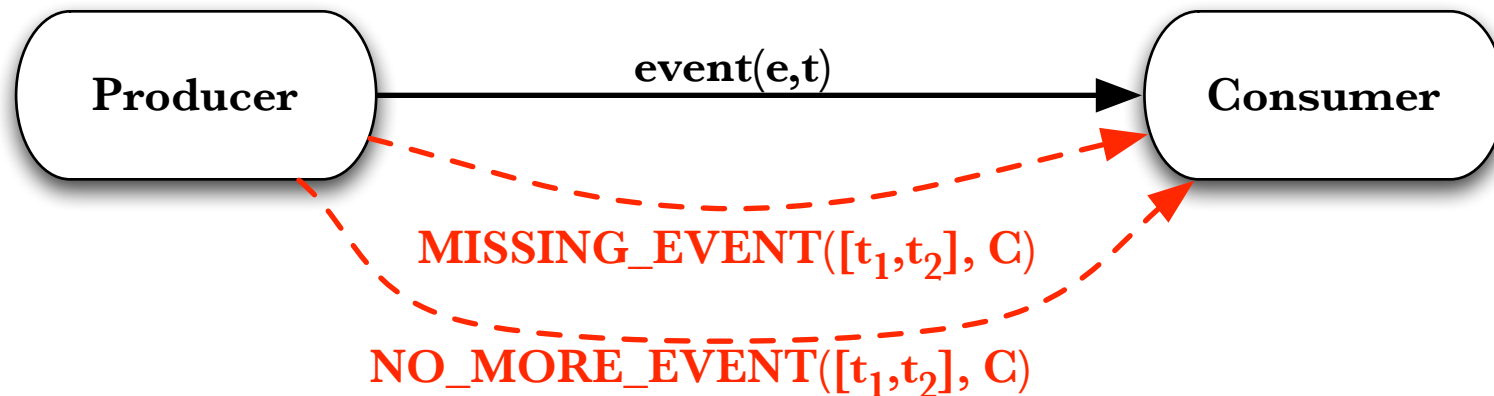


# Missed Observations

## Control message: **MISSING\_EVENT**

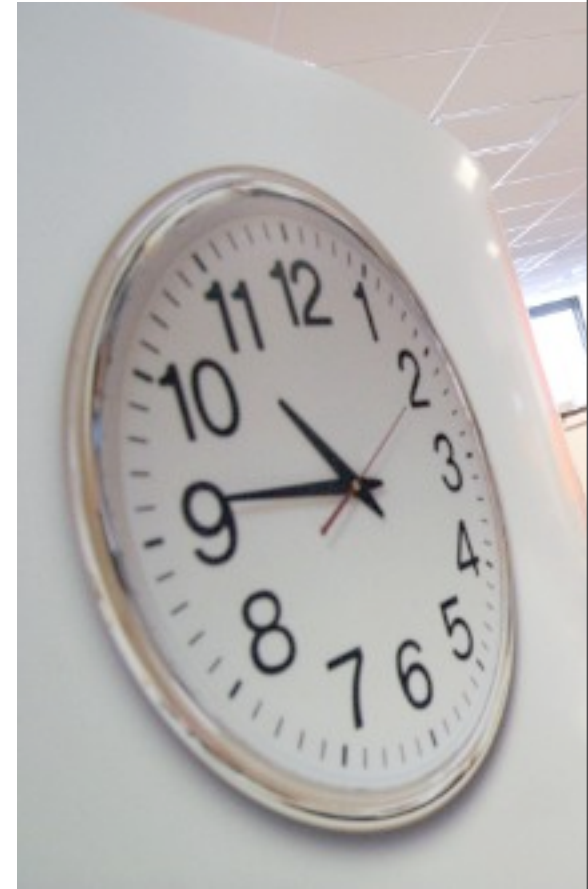
 *meaning: the producer indicates that some events may be missed (according to the specified condition and time window)*

- It could be due to a sensor pb (uncertainty) or to a lost of events (full buffer) or to a action of the producer (for algorithmic reasons, ...)



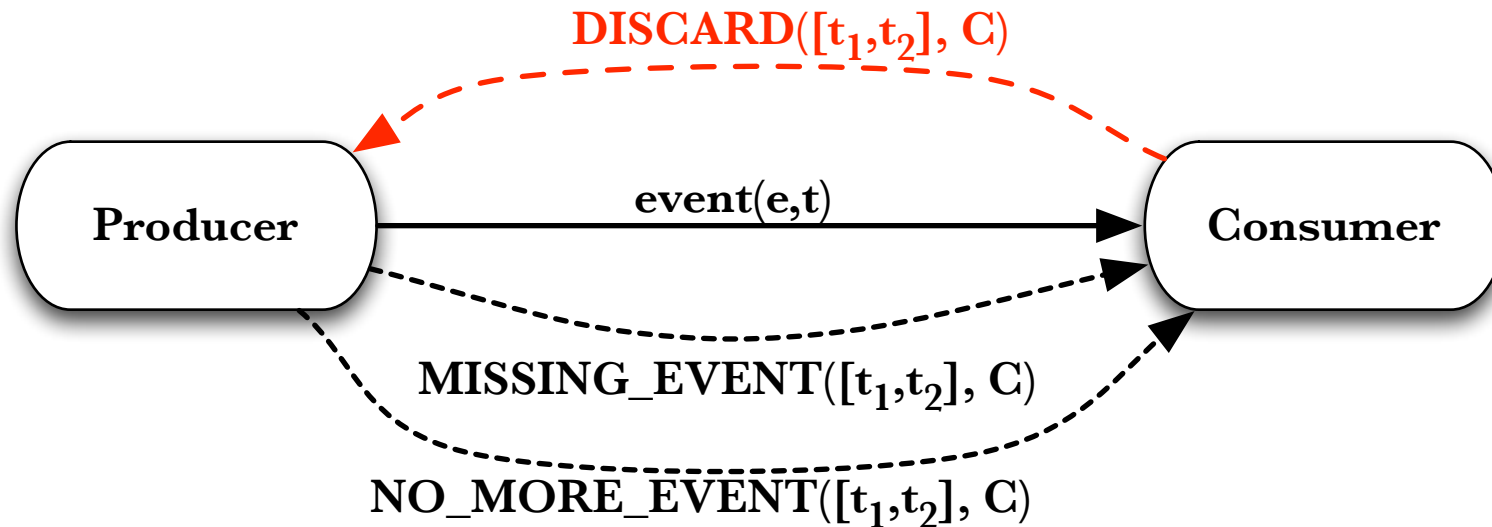
# 2

## Control by the Consumer *(feedback)*



# Obsolete Events

The consumer indicates that it does not need some events



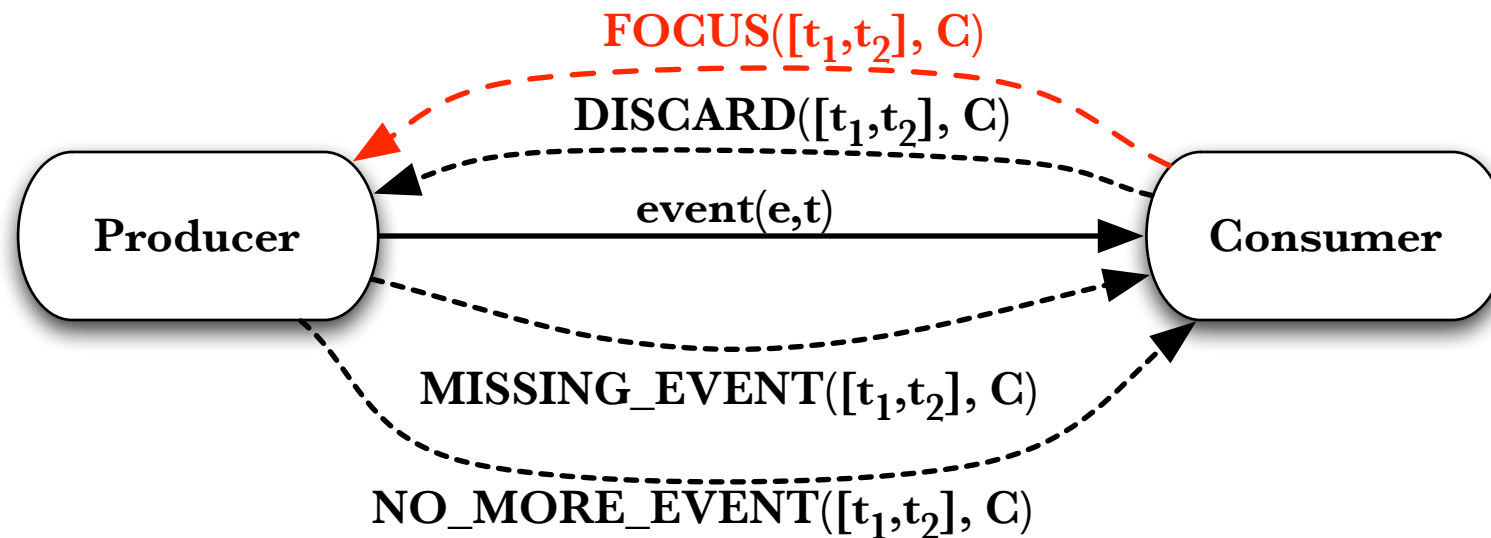
- *Note:* the producer could always send the events (i.e. ignore the DISCARD)
- *A contrario*, if the producer filters the events, it could inform the consumer with an appropriate NO\_MORE\_EVENT
- If the consumer retracts (REMOVE\_DISCARD), a MISSING\_EVENT message could be sent by the producer.



# Focusing

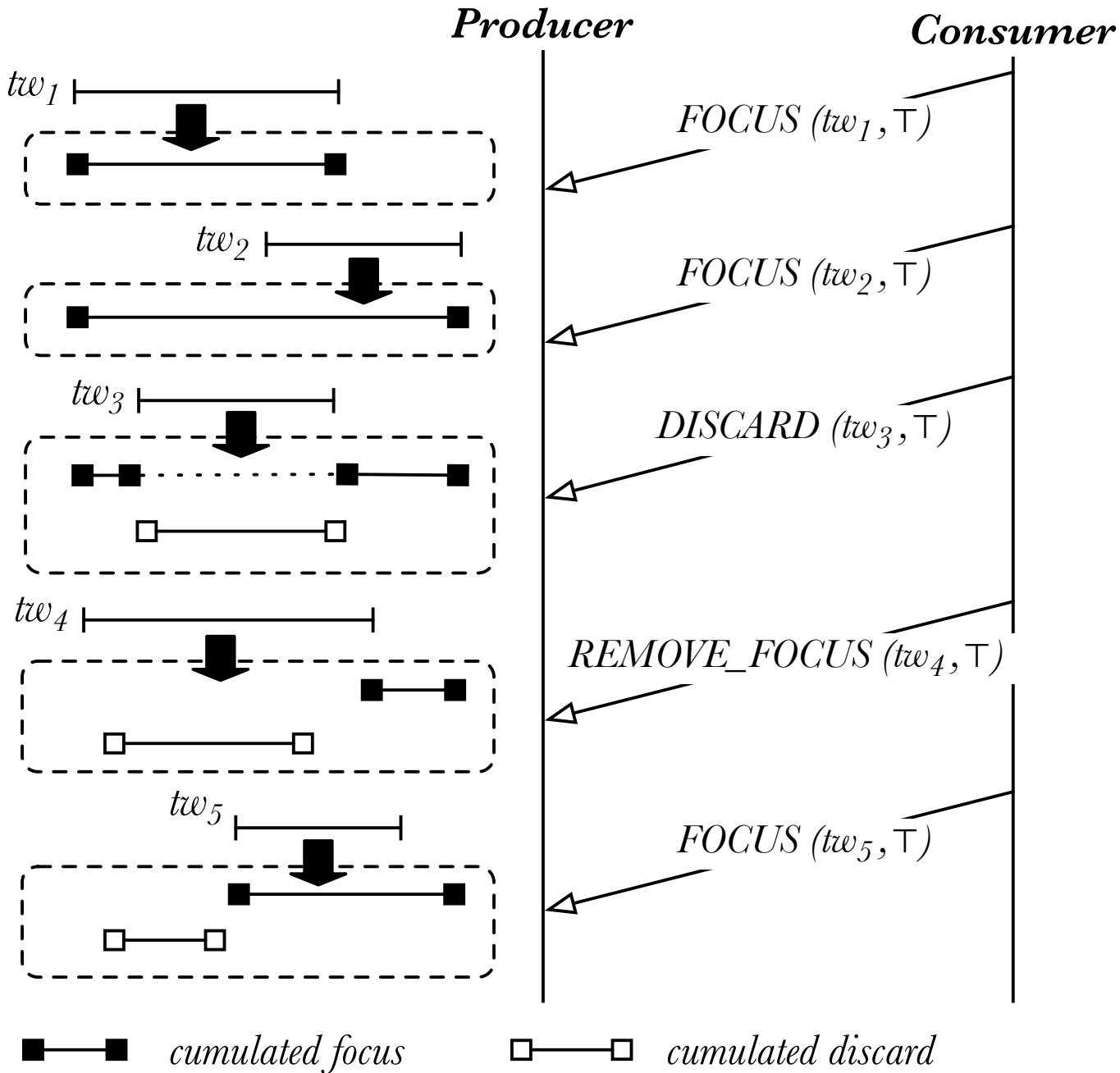
The consumer indicates that some events are “urgent”

 *contract: the producer should send the events as soon as possible*



- A REMOVE\_FOCUS message could cancel a focus
- Note: by default, a producer could delay its events as it desired if (and only if) it does not receive a FOCUS.

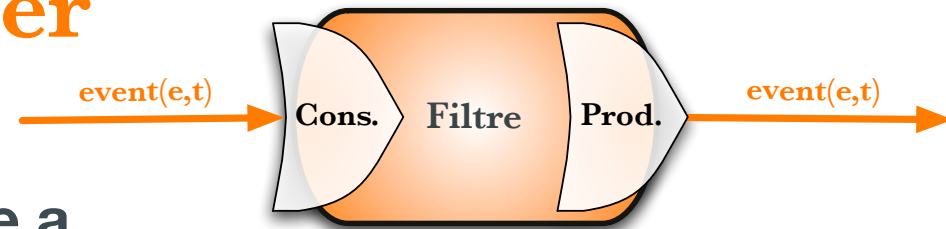




# Behavior when Cumulating FOCUS & DISCARD



# Example: impl. of a filter



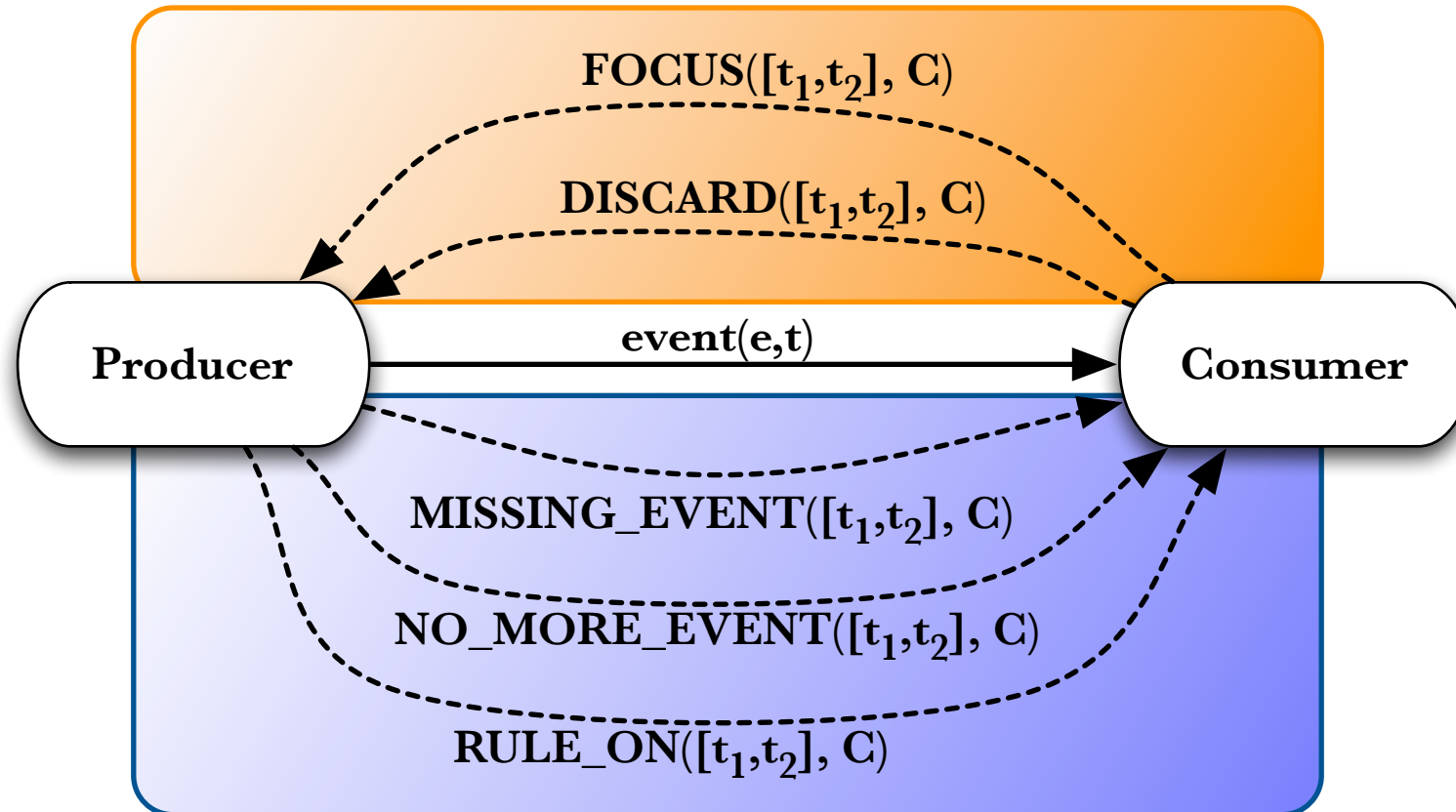
**This is a component which compute a filtering window (union of received DISCARD) and removes all the events which match this time window.**

- when a filter received a DISCARD(TW):
  - update:  $FILTER \leftarrow FILTER \cup TW$
  - send a NO\_MORE\_EVENT(TW) downstream and DISCARD(TW) upstream
- when a filter received a event (e,t):
  - if  $t \notin FILTER$ , send (e,t) downstream (if not, do nothing)
- when a filter received NO\_MORE\_EVENT(TW):
  - if  $TW \not\subset FILTER$ , send NO\_MORE\_EVENT(TW) downstream
- when a filter received FOCUS(TW):
  - if  $TW \not\subset FILTER$ , send FOCUS( $TW \cap FILTER$ ) upstream



# Summary

All control messages are  $\langle \text{TYPE} \rangle$  (TimeWindow, Condition)



- Note : the *REMOVE\_DISCARD* and *REMOVE FOCUS* messages are not represented here (but they still exist)



3

*Chronicles*



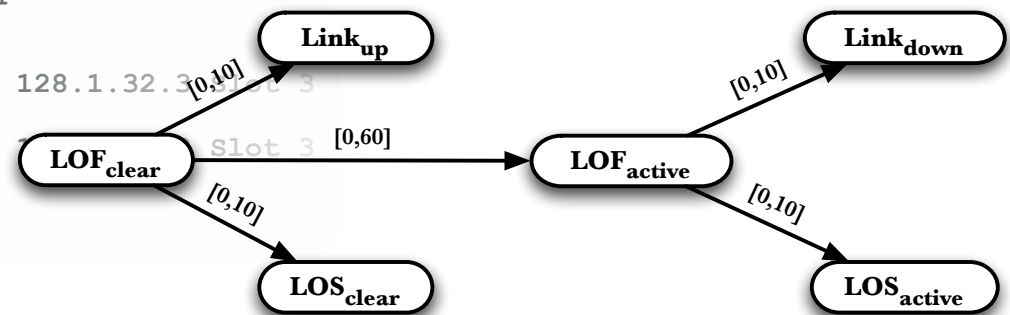


# Temporal Correlation

```
21/09/2001 17:36:32 lyon-2 Slot 2 Link 1 has gone down
21/09/2001 17:36:33 clermont Slot 1 Link 0 Path Yellow Alarm Active
21/09/2001 17:36:33 clermont Slot 1 Link 0 Path Label Mismatch Alarm Active
21/09/2001 17:36:34 clermont Slot 1 Link 0 has gone down
21/09/2001 17:36:34 grenoble Slot 5 Link 0 Path Yellow Alarm Active
21/09/2001 17:36:34 svcQsaalUpDown trap received from 128.1.16.2 Slot 2
21/09/2001 17:36:35 grenoble Slot 5 Link 0 has gone down
21/09/2001 17:36:36 grenoble Slot 5 Link 0 Path AIS Alarm Active
21/09/2001 17:36:36 svcQsaalUpDown trap received from 128.1.29.5 Slot 5
21/09/2001 17:37:42 sta-4 Slot 11 Link 0 Loss of Frame Alarm Clear
21/09/2001 17:37:43 sta-4 Slot 11 Link 0 Loss of Signal Alarm Clear
21/09/2001 17:37:51 svcQsaalUpDown trap received from 128.1.26.1 Slot 1
21/09/2001 17:37:55 svcQsaalUpDown trap received from 128.1.4.5 Slot 5
21/09/2001 17:37:57 sta-4 Slot 11 Link 0 has come up
21/09/2001 17:38:05 sta-4 Slot 11 Link 0 Loss of Frame Alarm Active
21/09/2001 17:38:26 sta-4 Slot 11 Link 0 has gone down
21/09/2001 17:38:36 grenoble Slot 5 Link 0 Path AIS Alarm Clear
21/09/2001 17:38:52 sta-4 Slot 11 Link 0 Loss of Signal Alarm Active
21/09/2001 17:42:11 lyon-2 Slot 1 has gone down
21/09/2001 17:42:30 lyon-2 Slot 4 has gone down
21/09/2001 17:42:30 lyon-2 Slot 3 has gone down
21/09/2001 17:42:39 svcQsaalUpDown trap received from 128.1.4.5 Slot 5
21/09/2001 17:42:43 grenoble Slot 5 Link 0 Path Yellow Alarm Clear
21/09/2001 17:42:44 grenoble Slot 5 Link 0 has come up
21/09/2001 17:43:03 lyon-2 Slot 9 has gone down
21/09/2001 17:43:09 lyon-2 Slot 3 has come up
21/09/2001 17:43:09 svcQsaalUpDown trap received from 128.1.32.3 Slot 3
21/09/2001 17:43:20 lyon-2 Slot 0 has gone down
21/09/2001 17:43:21 svcQsaalUpDown trap received from
```

If receiving  
**LOFclear** and then  
**LOSclear** and  
**LinkUp** within 10  
seconds, and...

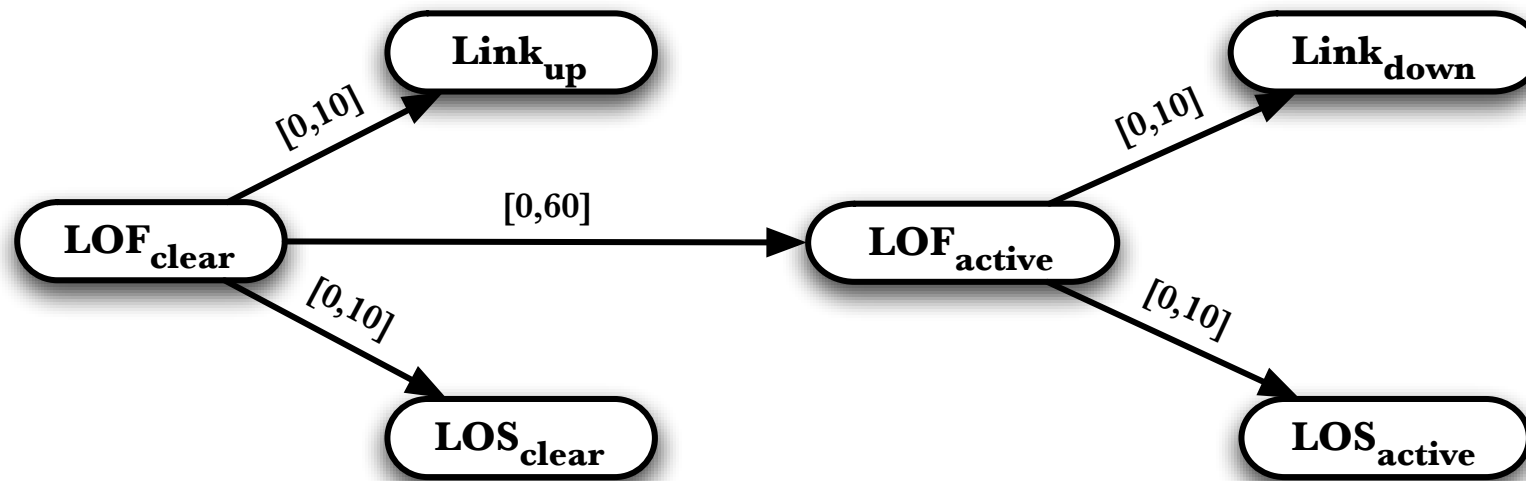
(we called a such  
temporal rule a  
*chronicle*)



# Temporal Correlation Rules

A such rule correlates a subset of events

- it is defined as a *set of patterns* of observable events...
- ... with *temporal constraints* between their *occurrence dates*
- example :



# Temporal Correlation Formalism (basics)

## Occurrence of event

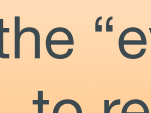
- `event ( M[196.3.2.4,?x,low] , t )`

  
the “event”  
to wait for

 occurrence date

## Absence of event

- `noevent ( M[102.0.0.1,?x,*] , (t1, t2) )`

  
the “event”  
to reject

  
lower  
bound

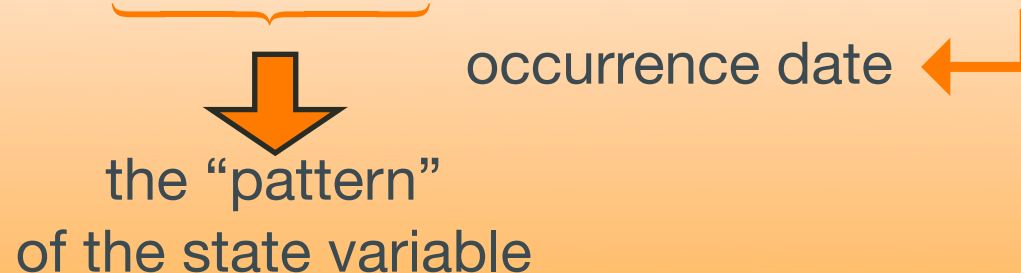
  
upper  
bound



# Temporal Correlation Formalism (more)

## Change of state variable

- `event( Severity[123,?x] : (low,high) , t )`



## Persistency of state

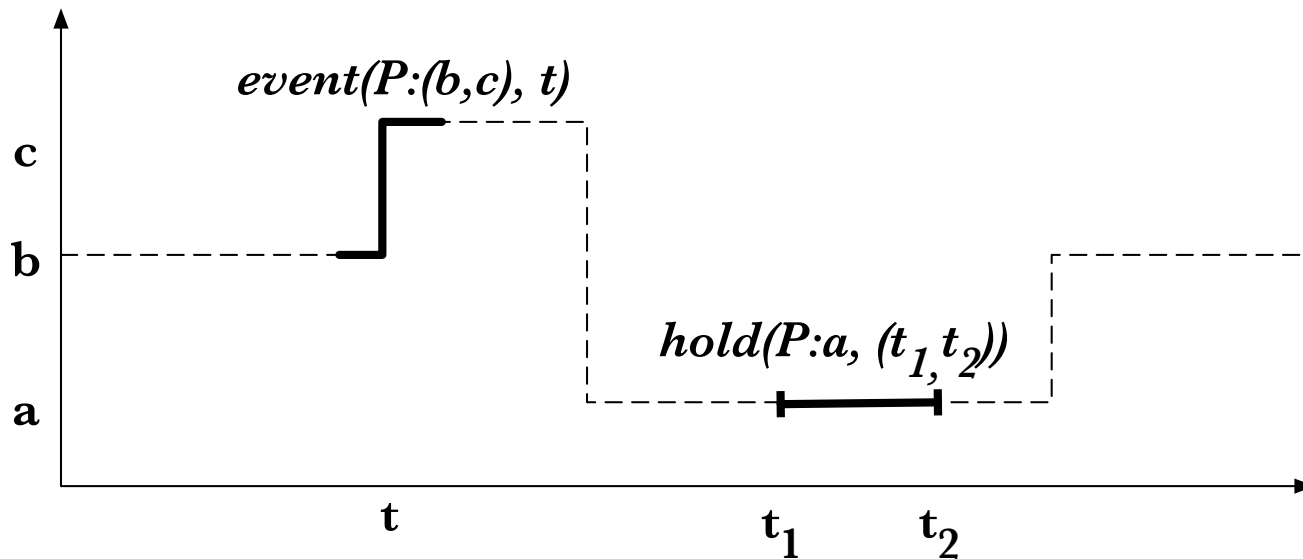
- `hold(Severity[123,?x] : low, (t1, t2) )`



# Temporal Correlation Formalism (more)

## Change of state variable

- `event( Severity[123,?x] : (low,high) , t)`



the “pattern”  
of the state variable

lower  
bound

upper  
bound



# An Example

Scenario : the **battery becomes critical** which is followed within 2 to 5 units of time by a loss of connectivity (good to bad) and a high level of IP packet loss less than 10 time units after the beginning. All of this scenario should be considered only for a WiFi connection.

```
chronicle DetectHandoverRequired_WIFI {  
    event(SOC: (? ,critical) ,t1) ←  
  
  
  
  
  
  
  
  
  
}
```



# An Example

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```
chronicle DetectHandoverRequired_WIFI {  
    event(SOC: (? ,critical) , t1)
```

```
    t2 - t1 in [2,5]
```



```
}
```



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```
chronicle DetectHandoverRequired_WIFI {  
    event(SOC: (? ,critical) , t1)  
    event(link: (good,bad) , t2) ←  
  
    t2 - t1 in [2,5]  
  
}
```





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```
chronicle DetectHandoverRequired_WIFI {  
    event(SOC: (? ,critical) , t1)  
    event(link: (good,bad) , t2)  
    event(IPloss: (? ,high) , t3) ←  
  
    t2 - t1 in [2,5]  
  
}
```



# An Example

Scenario : the battery becomes critical which is followed within 2 to 5 units of time by a loss of connectivity (good to bad) and a high level of IP packet loss **less than 10 time units after the beginning**. All of this scenario should be considered only for a WiFi connection.

```
chronicle DetectHandoverRequired_WIFI {  
    event(SOC: (? ,critical) , t1)  
    event(link: (good,bad) , t2)  
    event(IPloss: (? ,high) , t3)  
  
    t2 - t1 in [2,5]  
  
    t3 - t1 in [0,10] ←  
  
}
```



# An Example

Scenario : the battery becomes critical which is followed within 2 to 5 units of time by a loss of connectivity (good to bad) and a high level of IP packet loss less than 10 time units after the beginning.

**All of this scenario should be considered only for a WiFi connection.**

```
chronicle DetectHandoverRequired_WIFI {  
    event(SOC: (? ,critical) , t1)  
    event(link: (good,bad) , t2)  
    event(IPloss: (? ,high) , t3)  
    hold(connection:WiFi, (t1, t4)) ←  
    t2 - t1 in [2,5]  
    t2 < t4 ←  
    t3 - t1 in [0,10]  
    t3 < t4 ←  
}
```



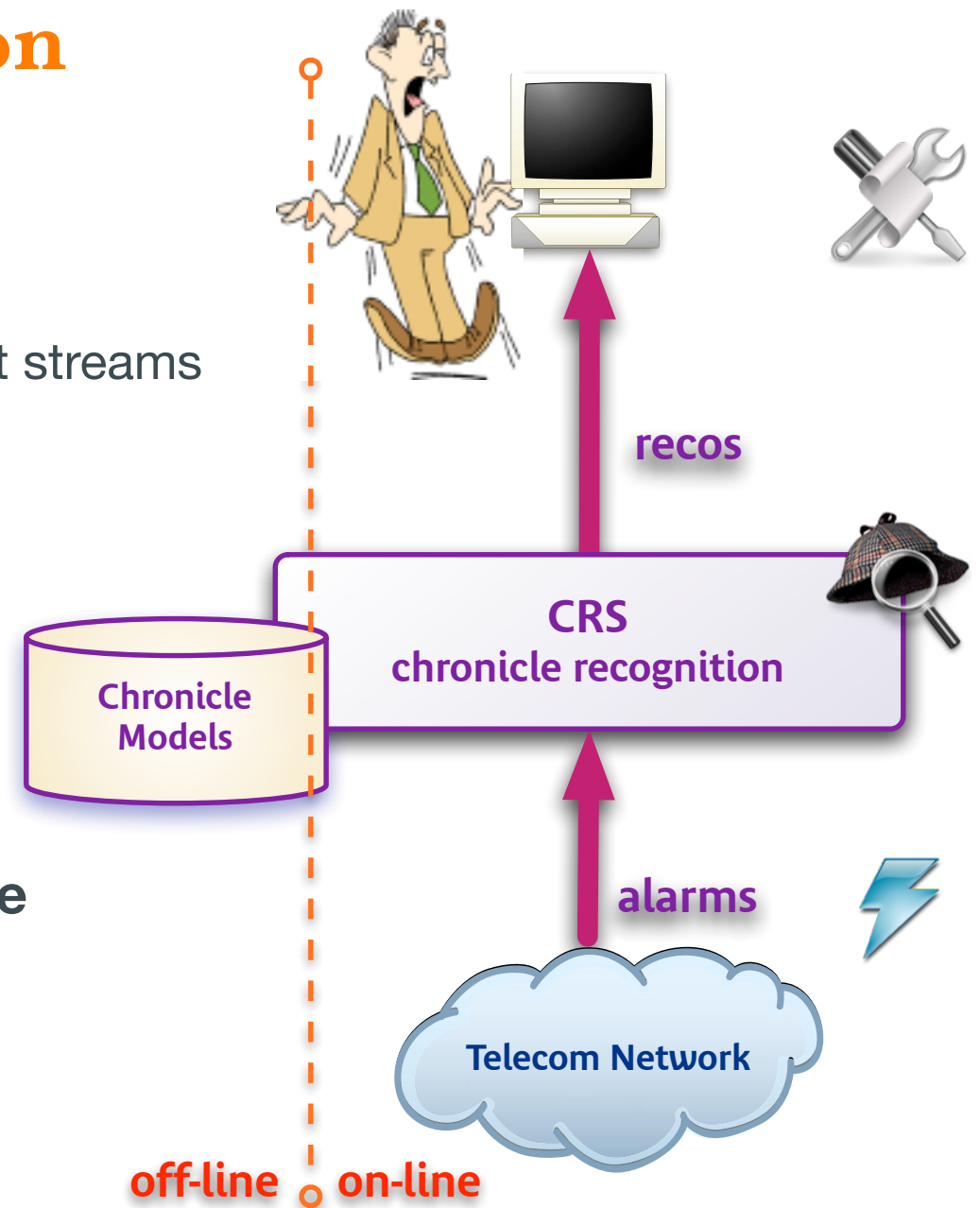
# Chronicle Recognition

## On-line Processing

- works on non-chronological event streams
- on the fly (one pass)
- 24h/24
- thousands evts/second

## chronicle language is expressive

- event, noevent, time constraints
- non necessary complete model



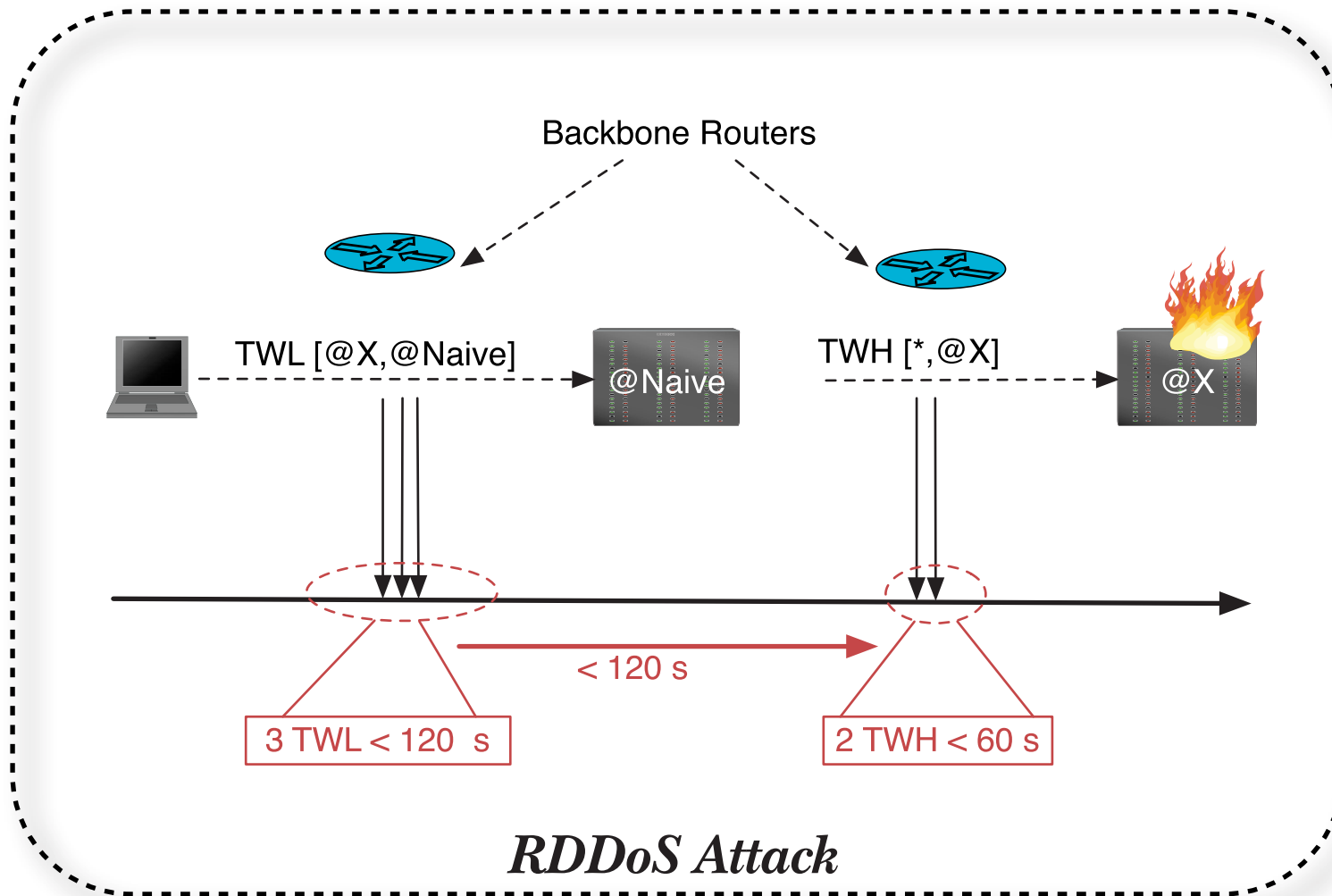
# 4

Usage for  
Correlation

*(example)*



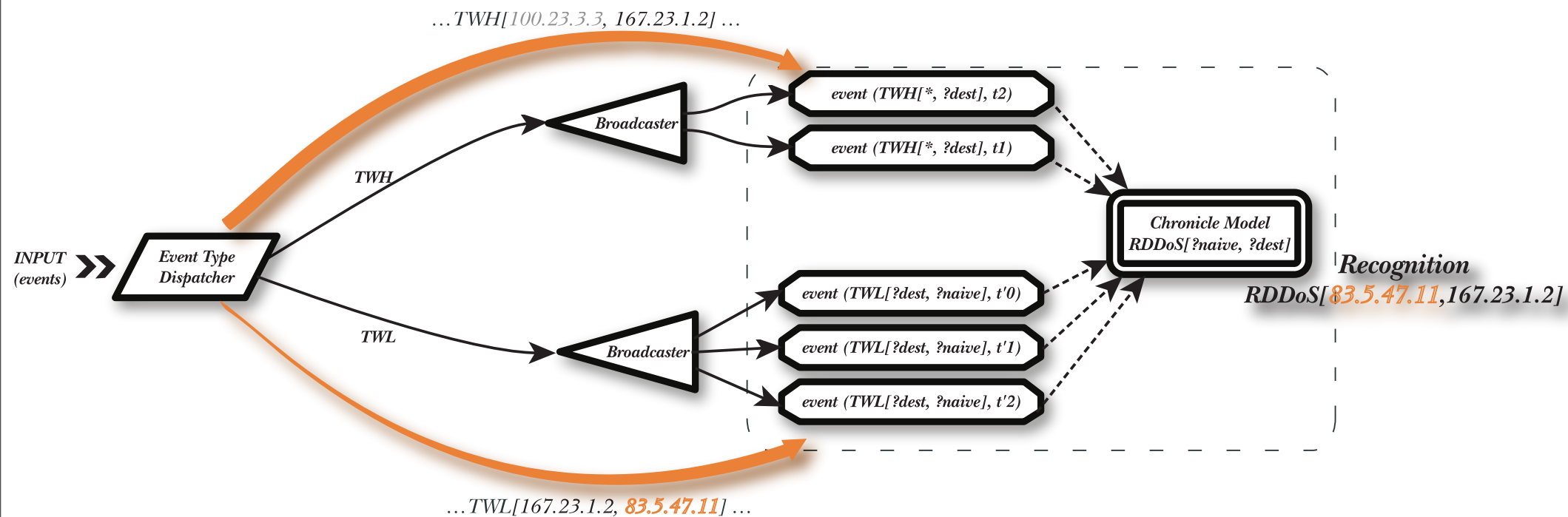
# Attack detection (1/4)



# Attack detection (2/4)

## Classical behavior of correlation system

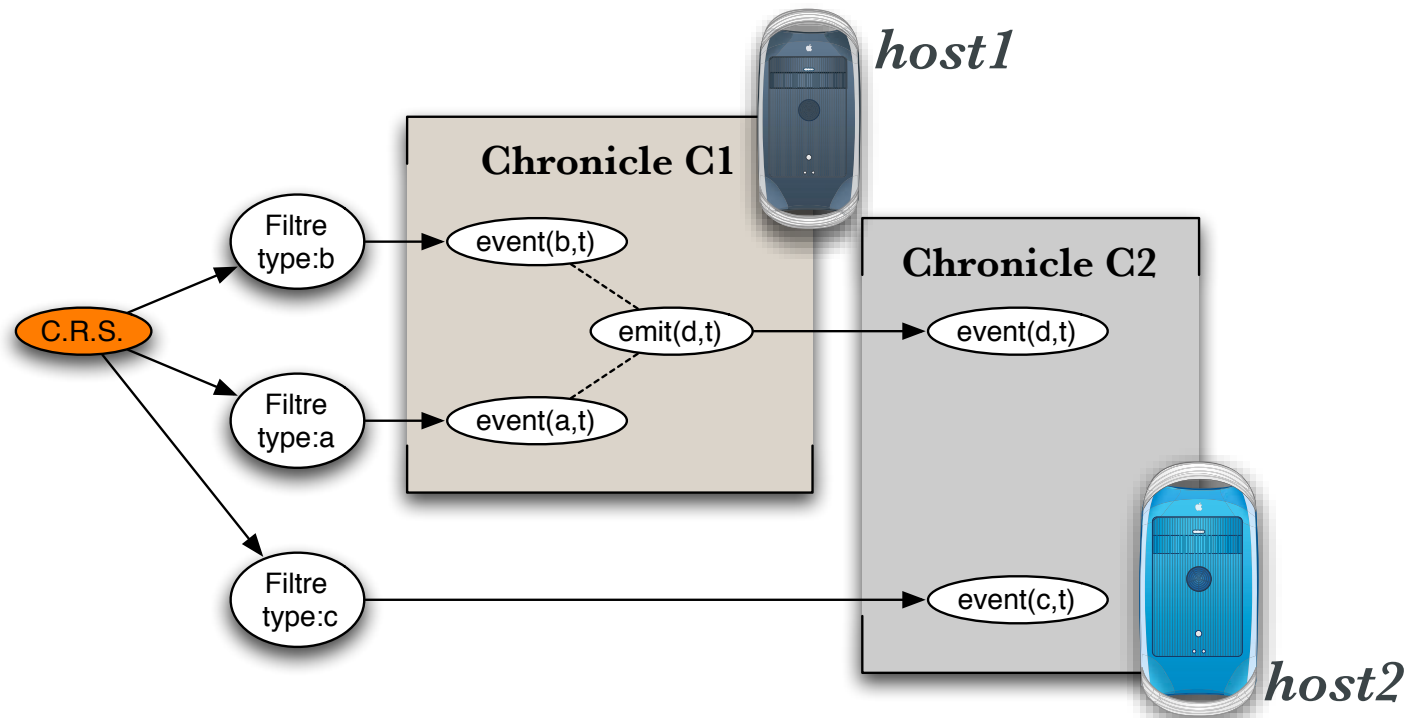
- events are integrated in their arrival order



# Distributed processing

## Principle

- Some temporal rules (chronicles) generates new event when recognized
- These chronicles are considered (and implemented) as producers





# Conclusion

## Modular architecture

- Addition of new components is easy: the processing is not only in the “code” but also in the structure (which could be automatically optimized and/or distributed)
- Impact on correlation tool (s.t. CRS): this is not a black box but a set of connected components → easy to distribute but also easy to modify the behavior (e.g. focusing)



## Currently on progress...

- Hot-connection of components
- Multi-hosts components (at the moment, just RMI components available)
- On editing and debugging tools (graphical editing)
- On optimisation tool (of the structure)



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Thank  
You!

Any question ?